

# Meanings as proposals: an algebraic semantics

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LOFT 2012, June 20<sup>th</sup>

# Goal

A semantics that distinguishes:

(1)  $p \vee q$

(2)  $p \vee q \vee (p \wedge q)$

# Approach

(Following Roelofsen 2011)

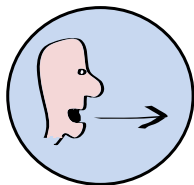
1. Choose a particular perspective on meaning.
2. Derive a formal semantics from this perspective.

## Example: deriving classical semantics

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## Meaning as information

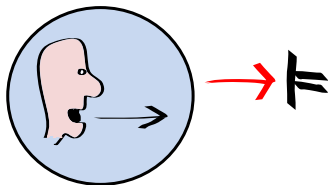
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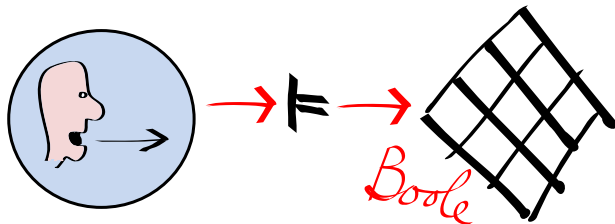
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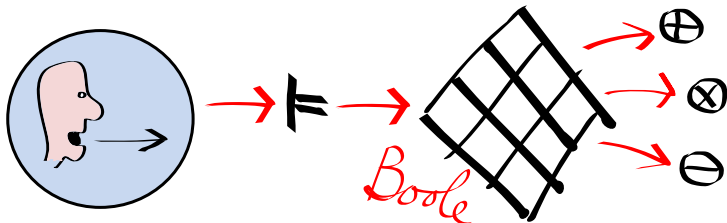
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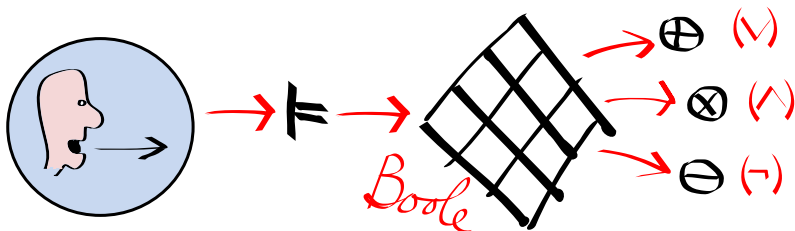




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## Deriving an algebra of proposals (1/2)

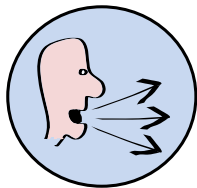


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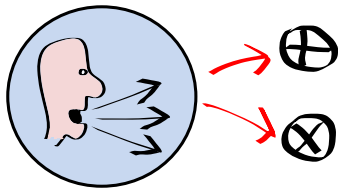


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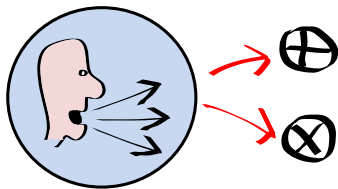
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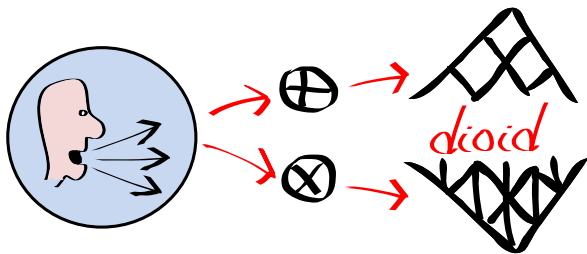
## Definition

$$A \otimes B := \{f \circ g : f \in A, g \in B\}$$

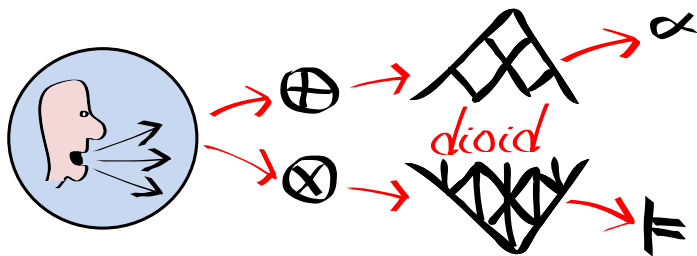
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- ▶  $\cup \widehat{[\varphi]} = [\varphi]_{\text{Classical}}$
  - ▶  $\widehat{[\varphi]} \downarrow = [\varphi]_{\text{Inquisitive}}$
- (Groenendijk, Roelofsen, et al.)

Fin.

Thanks to the Netherlands Organization for Scientific Research (NWO) for financial support; to F. Roelofsen, J. Groenendijk, J. Marti, I. Ciardelli and an anonymous reviewer for valuable comments.

**Miniature posters  
available!**

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